

# Block Ciphers and DES

## Examples

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### 1 Simplified DES Example

Assume input 10-bit key,  $K$ , is: 1010000010

Then the steps for generating the two 8-bit round keys,  $K_1$  and  $K_2$ , are:

1. Rearrange  $K$  using  $P_{10}$ : 1000001100
2. Left shift by 1 position both the left and right halves: 00001 11000
3. Rearrange the halves with  $P_8$  to produce  $K_1$ : 10100100
4. Left shift by 2 positions the left and right halves: 00100 00011
5. Rearrange the halves with  $P_8$  to produce  $K_2$ : 01000011

$K_1$  and  $K_2$  are used as inputs in the encryption and decryption stages.

Assume a 8-bit plaintext,  $P$ : 01110010

Then the steps for encryption are:

1. Apply the initial permutation,  $IP$ , on  $P$ : 10101001
2. Assume the input from step 1 is in two halves,  $L$  and  $R$ :  $L=1010$ ,  $R=1001$
3. Expand and permute  $R$  using  $E/P$ : 11000011
4. XOR input from step 3 with  $K_1$ :  $10100100 \text{ XOR } 11000011 = 01100111$
5. Input left half of step 4 into S-Box  $S_0$  and right half into S-Box  $S_1$ :
  - a. For  $S_0$ : 0110 as input:  $b_1, b_4$  for row,  $b_2, b_3$  for column
  - b. Row 00, column 11  $\rightarrow$  output is 10
  - c. For  $S_1$ : 0111 as input:
  - d. Row 01, column 11  $\rightarrow$  output is 11
6. Rearrange outputs from step 5 (1011) using  $P_4$ : 0111
7. XOR output from step 6 with  $L$  from step 2:  $0111 \text{ XOR } 1010 = 1101$
8. Now we have the output of step 7 as the left half and the original  $R$  as the right half. Switch the halves and move to round 2: 1001 1101
9.  $E/P$  with right half:  $E/P(1101) = 11101011$
10. XOR output of step 9 with  $K_2$ :  $11101011 \text{ XOR } 01000011 = 10101000$
11. Input to s-boxes:
  - a. For  $S_0$ , 1010
  - b. Row 10, column 01  $\rightarrow$  output is 10

- c. For S1, 1000
  - d. Row 10, column 00 -> output is 11
12. Rearrange output from step 11 (1011) using P4: 0111
  13. XOR output of step 12 with left halve from step 8: 0111 XOR 1001 = 1110
  14. Input output from step 13 and right halve from step 8 into inverse IP
    - a. Input us 1110 1101
    - b. Output is: 01110111

So our encrypted result of plaintext 01110010 with key 1010000010 is: 01110111

Other examples (encrypt or decrypt) could be:

- Plaintext: 11010101; Key: 0111010001; Ciphertext: 01110011
- Plaintext: 01001100; Key: 1111111111; Ciphertext: 00100010
- Plaintext: 00000000; Key: 0000000000; Ciphertext: 11110000
- Plaintext: 11111111; Key: 1111111111; Ciphertext: 00001111